

(-1)-Enumerations of arrowed Gelfand–Tsetlin patterns

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joint work with I. Fischer

Overview

- Arrowed Gelfand–Tsetlin patterns
- A (-1) -enumeration of arrowed GT patterns
- Proof Sketch

Classical GT pattern

A **Gelfand-Tsetlin pattern** (GT) is a triangular array of integers of the form

$$\begin{array}{c} T_{1,1} \\ T_{2,1} \quad T_{2,2} \\ \dots \quad \dots \\ T_{n,1} \quad \dots \quad T_{n,n} \end{array} \qquad \begin{array}{c} T_{i,j} \\ \swarrow \quad \searrow \\ T_{i+1,j} \leq T_{i+1,j+1} \end{array}$$

The weight of a GT pattern T is $\mathbf{x}^T := \prod_{i=1}^n x_i^{\sum_{j=1}^i (T_{i,j}) - \sum_{j=1}^{i-1} (T_{i-1,j})}$.

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For a partition $\lambda = (\lambda_1, \dots, \lambda_n)$, the **Schur polynomial** s_λ is

$$s_\lambda(\mathbf{x}) = \sum_T \mathbf{x}^T,$$

where the sum is over all GTs T with bottom row $\lambda_n, \lambda_{n-1}, \dots, \lambda_1$.

An Example

For $\lambda = (2, 2, 1)$ we have

$$\begin{array}{ccc} \begin{array}{c} 1 \\ 1 \ 2 \\ 1 \ 2 \ 2 \\ x_1 x_2^2 x_3^2 \end{array} & \begin{array}{c} 2 \\ 1 \ 2 \\ 1 \ 2 \ 2 \\ x_1^2 x_2 x_3^2 \end{array} & \begin{array}{c} 2 \\ 2 \ 2 \\ 1 \ 2 \ 2 \\ x_1^2 x_2^2 x_3 \end{array} \end{array}$$

$$s_{(2,2,1)}(x_1, x_2, x_3) = x_1 x_2^2 x_3^2 + x_1^2 x_2 x_3^2 + x_1^2 x_2^2 x_3.$$

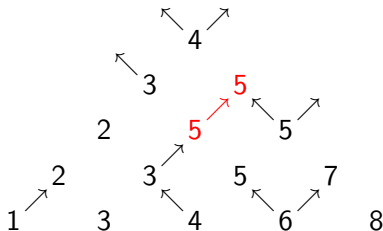
Arrowed Gelfand Tsetlin pattern

An **arrowed Gelfand-Tsetlin pattern** is a GT pattern $(T_{i,j})$ together with a decoration of the entries by the symbols $\emptyset, \swarrow, \nearrow, \nwarrow \nearrow$ such that

$$T_{i+1,j} = T_{i,j} \text{ and } T_{i+1,j} \text{ is decorated by } \nearrow \text{ or } \nwarrow \nearrow,$$

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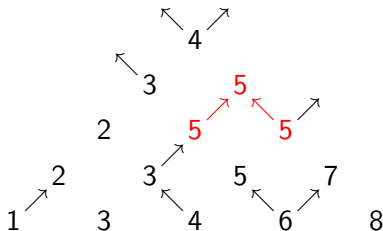
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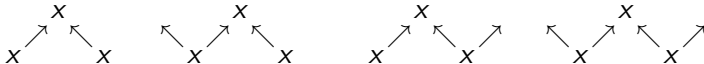
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The weight of an AGT

We call the following local configurations **special little triangles**



The sign of an AGT T is

$$\text{sgn}(T) = (-1)^{\# \text{ of special little triangles in } T}.$$

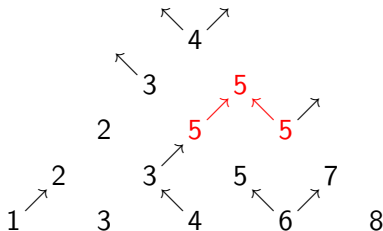
We define the weight $W(A)$ of A as

$$\text{sgn}(T) \cdot t^{\# \nearrow} u^{\# \searrow} v^{\# \swarrow} w^{\# \nwarrow} \cdot \mathbf{x}^T \prod_{i=1}^n x_i^{\# \nearrow \text{ in row } i - \# \searrow \text{ in row } i}.$$

An example

$$\text{sgn}(T) \cdot t^{\#\emptyset} u^{\#\nearrow} v^{\#\nwarrow} w^{\#\nearrow\searrow} \cdot \mathbf{x}^T \prod_{i=1}^n x_i^{\#\nearrow \text{ in row } i - \#\nwarrow \text{ in row } i}$$

The arrowed Gelfand–Tsetlin pattern



has weight $-t^7 u^3 v^2 w^3 x_1^4 x_2^3 x_3^5 x_4^6 x_5^5$.

A multivariate generating function for AGTs

Denote by E_x the *shift operator* $E_x f(x) = f(x + 1)$.

Theorem (Fischer – S.A., 2023)

The weighted enumeration $\mathcal{A}_\lambda(t, u, v, w; \mathbf{x})$ of all AGTs with bottom row $(\lambda_n, \lambda_{n-1}, \dots, \lambda_1)$ is given by

$$\begin{aligned} \mathcal{A}_\lambda(t, u, v, w; \mathbf{x}) &= \prod_{i=1}^n (ux_i + vx_i^{-1} + w + t) \\ &\times \prod_{1 \leq i < j \leq n} \left(t \text{id} + uE_{\lambda_j} + vE_{\lambda_i}^{-1} + wE_{\lambda_j}E_{\lambda_i}^{-1} \right) s_\lambda(\mathbf{x}). \end{aligned}$$

Known specialisations

$$t \# \emptyset \quad u \# \nearrow \quad v \# \nwarrow \quad w \# \nearrow \nwarrow$$

- $\mathcal{A}_\lambda(1, 0, 0, 0; \mathbf{x}) = s_\lambda(\mathbf{x}),$
- $\mathcal{A}_\lambda(0, 0, 1, 0; \mathbf{x}) = s_{(\lambda_1-n, \lambda_2-n+1, \dots, \lambda_n-1)}(\mathbf{x}),$

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- $\mathcal{A}_{(n, n-1, \dots, 1)}(0, u, v, w; \mathbf{x})$ yields a weighted enumeration of alternating sign matrices,
- $\mathcal{A}_{(2n, 2n-2, \dots, 2)}(0, u, v, w; \mathbf{x})$ yields a weighted enumeration of vertically symmetric alternating sign matrices.

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For this talk we are interested in

- $\mathcal{A}_\lambda(1, 1, 1, -1; \mathbf{x})|_{x_i=1}$ and $\mathcal{A}_\lambda(1, 1, 1, 0; \mathbf{x})|_{x_i=1}$.

The main result

Theorem (Fischer – S.A.)

For positive integers n, m we have

$$\sum_{0 \leq \lambda_n < \dots < \lambda_1 \leq m} \mathcal{A}_\lambda(1, 1, 1, -1; \mathbf{1}) = 2^n \prod_{i=1}^n \frac{(m - n + 3i + 1)_{i-1} (m - n + i + 1)_i}{\left(\frac{m-n+i+2}{2}\right)_{i-1} (i)_i},$$

$$\sum_{0 \leq \lambda_n < \dots < \lambda_1 \leq m} \mathcal{A}_\lambda(1, 1, 1, 0; \mathbf{1}) = 3^{\binom{n+1}{2}} \prod_{i=1}^n \frac{(2n + m + 2 - 3i)_i}{(i)_i},$$

where $(a)_i = (a)(a+1)\cdots(a+i-1)$ is the *Pochhammer symbol*.

The case $m = n - 1$

Setting $m = n - 1$ implies $\lambda = (n - 1, n - 2, \dots, 1, 0)$ and hence

$$\begin{aligned} 2^{-n} \mathcal{A}_{(n-1, n-2, \dots, 1, 0)}(1, 1, 1, -1; \mathbf{1}) &= 2^{n(n-1)/2} \prod_{i=0}^{n-1} \frac{(4i+2)!}{(n+2i+1)!} \\ &= 1, 4, 60, 3328, 678912, \dots \end{aligned}$$

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- configurations of the 20 vertex model in a certain domain, and
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This was proved by Koutschan and extended in a recent preprint by Corteel, Huang and Krattenthaler.

Overview of the proof of the main results

- 1 Obtain a determinant
- 2 Guess a (partial) LU decomposition.
- 3 Proof the LU decomposition

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 - Actually we need to do a case distinction: m even/odd.
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- 3 Proof the LU decomposition
 - It “suffices” to prove a hypergeometric triple sum.
 - For this we use Mathematica implementations of Sister Celine’s algorithm and creative telescoping.

Ad 1) Littlewood identities

We have the operator formula for evaluating \mathcal{A}_λ

$$\begin{aligned} \mathcal{A}_\lambda(t, u, v, w; \mathbf{x}) &= \prod_{i=1}^n (ux_i + vx_i^{-1} + w + t) \\ &\times \prod_{1 \leq i < j \leq n} \left(t \operatorname{id} + uE_{\lambda_j} + vE_{\lambda_i}^{-1} + wE_{\lambda_j}E_{\lambda_i}^{-1} \right) s_\lambda(\mathbf{x}). \end{aligned}$$

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The classical Littlewood identity is

$$\sum_{\lambda} s_{\lambda}(\mathbf{x}) = \prod_{i=1}^n \frac{1}{1-x_i} \prod_{1 \leq i < j \leq n} \frac{1}{1-x_i x_j},$$

where $\mathbf{x} = (x_1, \dots, x_n)$.

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The classical **bounded** Littlewood identity is

$$\sum_{\lambda \subseteq (m^n)} s_\lambda(\mathbf{x}) = \prod_{i=1}^n \frac{1}{1-x_i} \prod_{1 \leq i < j \leq n} \frac{1}{1-x_i x_j} \cdot \frac{1}{x_j - x_i} \\ \times \det_{1 \leq i, j \leq n} \left(x_i^{j-1} \quad -x_i^{m+2n-j} \right),$$

where $\mathbf{x} = (x_1, \dots, x_n)$.

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In the generalised setting we have

$$\begin{aligned} \sum_{0 \leq \lambda_n < \dots < \lambda_1 \leq m} \mathcal{A}_\lambda(1, 1, 1, w; \mathbf{x}) &= \prod_{i=1}^n \frac{1}{1 - x_i} \prod_{1 \leq i < j \leq n} \frac{1}{1 - x_i x_j} \cdot \frac{1}{x_j - x_i} \\ &\times \det_{1 \leq i, j \leq n} \left(x_i^{j-1} f_j(x_i) - x_i^{m+2n-j} f_j(x_i^{-1}) \right) \prod_{i=1}^n (x_i^{-1} + 1 + w + x_i), \end{aligned}$$

where $\mathbf{x} = (x_1, \dots, x_n)$ and $f_j(x) = (1 + x)^{j-1} (1 + wx)^{n-j}$.

Ad 1) A simple determinant

By setting $m = 2l + 1$ and $x_1 = \dots = x_n = 1$, we obtain

$$(3 + w)^{n2^n} \det_{1 \leq i, j \leq n} \left(\sum_{p, q} w^{n-j-q} (-1)^j \times \binom{j-1}{p} \binom{n-j}{q} \binom{p-q-\ell+i-2}{2i-1} \right).$$

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For $w = -1$ this can be simplified by using the Chu-Vandermonde identity

$$2^{2n} \det_{1 \leq i, j \leq n} \left(\sum_p \binom{n-j}{p} \binom{\ell-p+i}{2i-j} \right).$$

Ad 2) Guessing the LU decomposition

Define $a_{i,j} = \sum_p \binom{n-j}{p} \binom{\ell-p+i}{2i-j}$ and

$$x_{i,j} = \begin{cases} (-1)^{i+1} \frac{\binom{j}{j}}{(2\ell - n + 3j + 2)_{j-1} (2\ell - n + i + 2)_j} \\ \quad \times \sum_t \left(2^{2i-4t-n} (\ell - n/2 + j/2 + t + 3/2)_{i-2t-1} \right) & i \leq j, \\ \quad \times \frac{(i-j-2t+1)_{2t} (i-2j+1)_{j-1-t}}{(1)_t (1)_{i-2t-1}} \\ 0 & \text{otherwise.} \end{cases}$$

Lemma

We have

$$\sum_{k=1}^n a_{i,k} x_{k,j} = \begin{cases} 1 & i = j, \\ 0 & i < j. \end{cases}$$

Ad 3) Next steps

- The expression in the sum can be simplified by using transformations for hypergeometric series (the package HYP by Krattenthaler was very useful for this!).
- However, the triple sum can not be evaluated by summation rules for hypergeometric series (as far as we are aware of).

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- The expression in the sum can be simplified by using transformations for hypergeometric series (the package HYP by Krattenthaler was very useful for this!).
- However, the triple sum can not be evaluated by summation rules for hypergeometric series (as far as we are aware of).
- We use two algorithms (Sister Celine's method, creative telescoping) which provide recursions for the triple sum and allow us to prove the Lemma.

Idea of Sister Celine's method

- Given a function $F(n) = \sum_k f(n, k)$ which we want to evaluate,
 - in our case: we want to show $F(n) = 0$ or $F(n) = 1$,
 - $f(n, k)$ consists of Pochhammer symbols.

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 - in our case: we want to show $F(n) = 0$ or $F(n) = 1$,
 - $f(n, k)$ consists of Pochhammer symbols.
- Assume we can find a recursion for f of the form

$$\sum_{r,s} a_{r,s}(n) f(n-r, k-s) = 0,$$

then we obtain

$$0 = \sum_k \left(\sum_{r,s} a_{r,s}(n) f(n-r, k-s) \right) = \sum_{r,s} a_{r,s} F(n-r).$$

Basic idea of creative telescoping

- Given a function $F(n) = \sum_{k=i}^j f(n, k)$ which we want to evaluate.
- Assume we can find a recursion

$$a(n)f(n, k) + b(n)f(n+1, k) = g(n, k+1) - g(n, k),$$

- then we obtain for $F(n)$

$$a(n)F(n) + b(n)F(n+1) = g(n, j+1) - g(n, i).$$

Careful checking is necessary!

Let $f(n) = (n+1)_n$ and remember

$$(x)_n = \begin{cases} (x)(x+1)\cdots(x+n-1) & n > 0, \\ 1 & n = 0, \\ \frac{1}{(x-1)(x-2)\cdots(x+n)} & n < 0. \end{cases}$$

The above algorithms will yield the recursion

$$f(n) = 2(2n-1)f(n-1),$$

which is however only true if $n \neq 0$.

